

# Mathematics of Infrastructure Planning

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## Exercise sheet 9

Deadline: Thu, Jun 14, 2012, **23:59**, <mailto:borndorfer@zib.de>

### Exercise 30.

10 points

Consider the following discrete stop location problems in a network  $N = (S \cup T, E)$  with demand points  $V$  and covering radius  $r$ :

- (DSL)  $p = |U|/S / \text{cov}_r(U) = V/\ell_2/p$  (planning from scratch)
- (DSL1)  $p = |U|/T / \text{cov}_r(U) = V/\ell_2/p$  (closing stops)
- (DSL2)  $p = |U|/S / \text{cov}_r(U \cup T) = V/\ell_2/p$  (opening stops)
- (DSL3)  $p = |U|/S \cup T / \text{cov}_r(U) = V/\ell_2/p$  (closing and opening stops).

Prove that (DSL $_i$ ) can be reduced to (DSL),  $i = 1, 2, 3$ .

### Exercise 31.

10 points

Consider a set covering problem (SCP)  $\min c^T x, Ax = \mathbf{1}, x \in \{0, 1\}^n$  with constraint matrix  $A \in \{0, 1\}^{m \times n}$  and objective  $c \in \mathbb{R}_+^n$ . Prove the validity of the following preprocessing rules:

- a)  $A_i = e_j \implies x_j = 1$  in every solution of (SCP).
- b)  $A_j \leq A_k$  and  $c_j < c_k \implies x_j = 0$  in every optimal solution of (SCP).
- c)  $A_i \leq A_k \implies A_k x \geq 1$  is redundant.
- d) Find, formulate, and prove another preprocessing rule.

### Exercise 32.

10 points

Let  $w^1 \geq \dots \geq w^s > w^{s+1} = 0$  be integers and denote by  $H(n) := \sum_{i=1}^n \frac{1}{i}$  the  $n$ -th harmonic number,  $n \in \mathbb{N}$ . Prove that

$$\sum_{i=1}^s (w^i - w^{i+1})/w^i \leq \sum_{i=1}^s (H(w^i) - H(w^{i+1})) \leq H(w^1).$$

### Exercise 33.

(Tutorial session)

Consider the bus stop location problem on the left of Fig. 1 with unit costs, potential stops  $S = \{s_1, \dots, s_7\}$ , demand points  $V = \{v_1, \dots, v_5\}$  and covering radius  $r$  with

associated covering matrix

$$A_r^{\text{cov}} = \left( \begin{array}{cccccc|c} s_1 & s_2 & s_3 & s_4 & s_5 & s_6 & s_7 & \\ \hline 1 & 0 & 0 & 0 & 0 & 0 & 0 & p_1 \\ 0 & 1 & 1 & 1 & 1 & 0 & 0 & p_2 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & p_3 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & p_4 \\ 1 & 1 & 1 & 1 & 0 & 1 & 0 & p_5 \end{array} \right)$$

(cf. Schöbel [2003]). Solve this problem using set covering techniques, illustrating your reasoning in the figure.

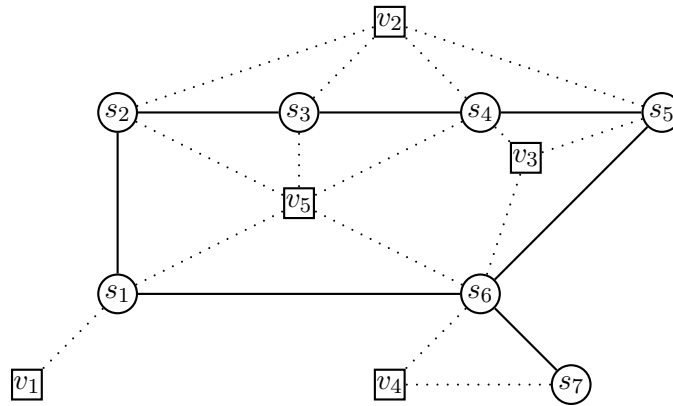


Figure 1: Bus stop location problem.

**Exercise 34.**

**(Tutorial session)**

Consider a *discrete stop location problem on a line*, i.e., in a network  $N = (S, E)$  of collinear nodes. Prove:

- The covering matrix  $A_r^{\text{cov}}$  associated with a discrete stop location problem on a line has the *consecutive ones property*, i.e., the columns and rows of  $A_r^{\text{cov}}$  can be permuted in such a way that  $a_{ri} = 1 = a_{rk}$  implies  $a_{rj} = 1, i \leq j \leq k$ , for all  $r \in V, i, k \in S$ .
- A 0/1 matrix with the consecutive ones property is totally unimodular.
- The set covering problem associated with a discrete stop location problem on a line can be solved by linear programming.

**Exercise 35.**

**(Tutorial session)**

Consider a continuous stop location problem in a network  $N = (S, E)$  with demand points  $V$  and covering radius  $r$ :

$$(\text{CSL}) \quad p = |U|/N / \text{cov}_r(U) = V/\ell_2/p.$$

Prove: If (CSL) is feasible, the optimum is finite.

**Exercise 36.****(Tutorial session)**

Consider a continuous stop location problem in a network  $N = (S, E)$  with demand points  $V$  and covering radius  $r$

$$\text{(CSL)} \quad p = |U|/N / \text{cov}_r(U) = V/\ell_2/p$$

and denote

$$S_r := S \cup \{x \in N : \exists v \in V : \|x - v\|_2 = r\}.$$

Prove that (CSL) can be reduced to (DSL):

- a)  $S_r$  is finite.
- b) Consider an edge  $e = uv \in E$  and let  $S_r^e := S_r \cap e = \{s_1, s_2, \dots, s_k\}$ . Order  $\{s_1, \dots, s_k\}$  as  $u = s_1 <_e \dots <_e s_k = v$  with respect to the natural order along the edge  $e = uv$ , starting from  $u$ . Let  $s$  be a point on  $uv$  between  $s_i$  and  $s_{i+1}$ , i.e.,  $s_i <_e s <_e s_{i+1}$ . Then  $\text{cov}_r(s_i) \supseteq \text{cov}_r(s)$  and  $\text{cov}_r(s_{i+1}) \supseteq \text{cov}_r(s)$ .
- c) (CSL) has an optimal solution  $S^* \subseteq S_r$ .