Advanced Practical Programming for Scientists Parallel Programming with OpenMP

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Sequential program

- From programmers perspective: Statements are executed in the order in which they appear
- On the CPU level this is not true
 - Instruction reordering due to compiler optimizations
 - Inside the CPU: Out of order execution to better utilize processing units

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Parallel program

- For the CPU completely independent programs executed on different cores
- Order of memory writes in one thread can be different when observed from another thread
- Programmer is responsible for all synchronization
- Statements from different threads can be interleaved in millions of ways even for small programs



What can happen if these statements are executed in parallel?

- Sequentially executed this program can only print "x is 1"
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Memory operations might not be visible in the same order when observed from the other thread!

Safe abstractions and memory model required for parallel programs



- OpenMP standard provides annotations for C, C++ , and Fortran languages
- If compiled without OpenMP the program is still a valid sequential program
- Compiler takes care of thread handling and provides abstractions for synchronization



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```
double sum = 0.0;
double vec[N];
//initialize vec
#pragma omp parallel for reduction(+:sum)
for ( int i = 0; i < N; ++i )
        x += vec[i];
```

Parallel section and work sharing



- #pragma omp parallel starts parallel section
- Parallel section is executed by all threads

```
#pragma omp parallel numthreads(4)
   implicit flush
   // executed by all 4 threads
   printf("thread_%i:_hello_world!\n",
        omp_get_thread_num());
   #pragma omp single
      // executed on one thread
   } // implicit barrier and flush
   // execute do_work1() and do_work2() in
        parallel
   #pragma omp sections
      #pragma omp section
      do_work1():
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      do_work2();
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- Work sharing constructs are used to distribute work
 - #pragma omp for
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 - #pragma omp single
- Memory consisitency enforced with #pragma omp flush operation
- Wait for threads to reach specified point with #pragma omp barrier

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Sharing state



- Default uses scoping rules to determines whether variables are shared
- Can also be specified explicitly with clauses
 - firstprivate(var)
 - private(var)
 - lastprivate(var)
 - shared(var)
- shared(var) clause only useful with default(none) or default(private) clause

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```
int x = 0; // x is shared (scope)
int y = 0; // y is private (clause)
#pragma omp parallel numthreads(4), firstprivate(y)
// private(y) would make y uninitialized
{
    int z = 0; // z is private (scope)
}
```

...



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- #pragma omp simd for instruction-level parallelism (since OpenMP 4.0)
- Directives for GPU computing (since OpenMP 4.0)

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Thank you for your attention!