Freie Universität Berlin FB Mathematik und Informatik Sommersemester 2017 Prof. Dr. Ralf Borndörfer Torsten Klug

Optimierung I

Exercise 1.

Consider the following LP

- a) Use the simplex algorithm to show that the optimal solution is $x_1^* = \frac{1}{2}, x_2^* = \frac{1}{2}$ and $x_3^* = 0$, by starting from the initial basic feasible solution: $x_1 = 0, x_2 = \frac{1}{2}, x_3 = \frac{1}{4}$.
- b) Write down the dual problem of the given LP.
- c) Use the optimal basic feasible solution from part a) to construct an optimal solution for the dual problem in part b).

Exercise 2.

- a) True or False: Any LP that has an optimal solution has either a unique optimal solution, or infinitely many optimal solutions. Reformulate the
- b) Give an example of an LP for which every feasible solution is optimal.
- c) True or False: If an LP is unbounded, then its dual problem is infeasible.

Exercise 3.

- a) Verify that $(x_1^*, x_2^*, x_3^*) = (0, 1, 1)$ is an optimal solution, and find the optimal dual values.
- b) It the optimal solution unique? Is it degenerated?
- c) Compute the range of all possible objective coefficients of x_3 for which $(x_1^*, x_2^*, x_3^*) = (0, 1, 1)$
- d) If the right-hand side of the first constraint decreases to -2, is x^* still optimal? What is the new optimal objective function value?

Exercise 4.

Let $x_1 \in \{0, 1\}, x_2 \in \{0, 1\}, x_5 \in \{0, 1\}, x_3 \ge 0, x_4 \ge 0.$

Reformulate the following constraints into linear constraints:

a)
$$5x_1x_2x_5 + 7x_3 + 8x_4 + |x_3 - 2x_4|$$

b)
$$\max\left\{\frac{4+6x_3}{1+x_1}, \frac{3x_1+5x_4}{2-x_2}\right\} \le 7$$

Exercise 5.

Let $c, a, b \in \mathbb{Z}^n$ be given and consider the equality knapsack problem:

$$\max c^T x$$
, s.t. $a^T x = b, x \ge 0$

- a) Derive conditions on c, a, and b under which the relaxation is feasible and bounded?
- b) Assuming that problem is both feasible and bounded, characterize its optimal solution(s).

Exercise 6.

A primal solution x and a dual solution y are each optimal for their respective problem if

- x is feasible for the primal,
- y is feasible for the dual, and
- x and y satisfy the complementary slackness.

How many and which ones of these three properties are satisfied at every iteration of the primal(i.e.,phase two) simplex algorithm.

Homepage of the Lecture: http://www.zib.de/ss17_Optimierung_I Questions?: klug@zib.de