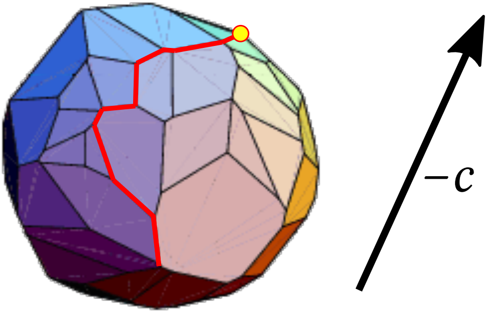
Introduction to

Linear and Combinatorial Optimization

4 Simplex Method

4.1 Basic Version

Basic Idea



Rough Description of Simplex Algorithm

- Start from a basic feasible solution
- In each iteration, move to a better adjacent vertex
- ... until no further improvement can be found

Linear Program in Standard Form

Throughout this section, we consider the following standard form problem:

$$\begin{aligned} \text{minimize} \quad & c^\top x \\ \text{subject to} \quad & A \cdot x = b \\ & x \geq 0 \end{aligned}$$

with $A \in \mathbb{R}^{m \times n}$, $\operatorname{rank}(A) = m$, $b \in \mathbb{R}^m$, and $c \in \mathbb{R}^n$

Some Notation

- a basis is a vector $B = (B(1), \dots, B(m))$ with

$$\{B(1), \dots, B(m)\} \subseteq \{1, \dots, n\} \text{ and } \operatorname{rank}(A_{B(1)}, \dots, A_{B(m)}) = m$$

- for a basis B, a corresponding non-basis is a vector

$$N = (N(1), \dots, N(n-m)) \text{ with}$$
$$\{B(1), \dots, B(m)\} \cup \{N(1), \dots, N(n-m)\} = \{1, \dots, n\}$$

- we write $j \in B$ if $j = B(i)$ for some $i \in \{1, \dots, m\}$ and $j \in N$ if $j = N(i)$ for some $j \in \{1, \dots, n-m\}$
- for $x \in \mathbb{R}^n$, the basic vector is $x_B = (x_{B(1)}, \dots, x_{B(m)})$ and the non-basic vector is $x_N = (x_{N(1)}, \dots, x_{N(n-m)})$
- the basic matrix is $A_B = [A_{B(1)}, \dots, A_{B(m)}]$ and
- the non-basic matrix is $A_N = [A_{N(1)}, \dots, A_{N(n-m)}]$

$$Ax = b \iff \sum_{j=1}^{n} A_j x_j = b \iff A_B x_B + A_N x_N = b$$

Basic Directions

Observation 4.1 The values of the basic variables x_B in the system $A \cdot x = b$ are uniquely determined by the values x_N of the non-basic variables.

Proof:

$$A \cdot x = b \iff A_B \cdot x_B + A_N x_N = b$$
$$\iff x_B = A_B^{-1} b - \sum_{j \in N} A_B^{-1} A_j x_j \quad \square$$

- for fixed $j \in N$, let $d \in \mathbb{R}^n$ be given by

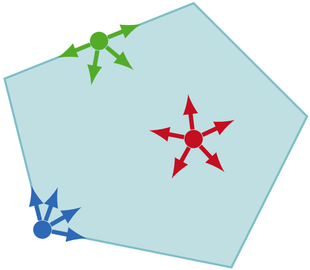
$$d_j := 1, \quad d_{j'} := 0 \text{ for } j' \in N \setminus \{j\} \text{ and } d_B := -A_B^{-1} \cdot A_j.$$

- then $A \cdot (x + \theta\, d) = b$, for all $\theta \in \mathbb{R}$
- d is called the jth basic direction

Feasible Directions

Definition 4.2 Let $P \subseteq \mathbb{R}^n$ a polyhedron. For $x \in P$ the vector $d \in \mathbb{R}^n \setminus \{0\}$ is a feasible direction at x if there is a $\theta > 0$ with $x + \theta\, d \in P$.

Example: Some feasible directions at several points of a polyhedron.



Basic Directions and Feasible Directions

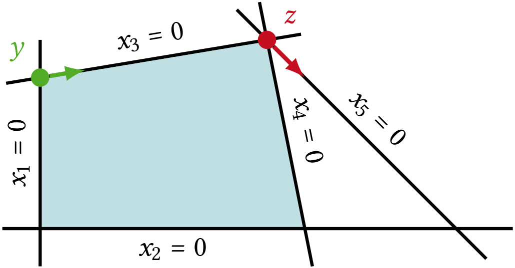
Consider a basic feasible solution x.

Question: Is the jth basic directions d a feasible direction?

Case 1: If x is a non-degenerate feasible solution, then $x_B > 0$ and $x + \theta\, d \geq 0$ for $\theta > 0$ small enough. $\longrightarrow$ answer is yes!

Case 2: If x is degenerate, the answer might be no!
E.g., if $x_{B(i)} = 0$ and $d_{B(i)} < 0$, then $x + \theta\, d \not\geq 0$, for all $\theta > 0$.

Example: $n = 5$, $m = 3$, $n - m = 2$



- 1st basic direction at y is feasible
  (basic variables $x_2,\ x_4,\ x_5$)
- 3rd basic direction at z is infeasible
  (basic variables $x_1,\ x_2,\ x_4$)

Reduced Cost Coefficients

Consider a basic solution x.

Question:

How does the cost change when moving along the jth basic direction d?

$$c^\top \cdot (x + \theta\, d) = c^\top \cdot x + \theta\, c^\top \cdot d = c^\top \cdot x + \theta\, \underbrace{(c_j - c_B{}^\top \cdot A_B^{-1} \cdot A_j)}_{\bar{c}_j\ :=}$$

Definition 4.3 For a given basis B and corresponding basic solution x, the reduced cost of variable x_j, $j = 1, \ldots, n$, is

$$\bar{c}_j := c_j - c_B^\top A_B^{-1} \cdot A_j.$$

Observation 4.4 The reduced cost of a basic variable $x_{B(i)}$ is zero.

Proof: $\bar{c}_{B(i)} = c_{B(i)} - c_B^\top \cdot \underbrace{A_B^{-1} \cdot A_{B(i)}}_{= e_i} = c_{B(i)} - c_{B(i)} = 0$ □

Optimality Criterion

Theorem 4.5 Let x be a basic feasible solution and $\bar{c}$ the vector of reduced costs.

a. If $\bar{c} \geq 0$, then x is an optimal solution.

b. If x is an optimal solution and non-degenerate, then $\bar{c} \geq 0$.

Definition 4.6 A basis B (or a basis matrix A_B) is optimal if

a. $A_B^{-1} \cdot b \geq 0$ and

b. $\bar{c}^\top = c^\top - c_B^\top \cdot A_B^{-1} \cdot A \geq 0$.

Observation 4.7 If B is an optimal basis, the associated basic solution x is feasible and optimal.

Proof of Theorem 4.5

a

Let B be the basis corresponding to x and let $y \in P$. Then,

$$\begin{aligned} c^\top \cdot y &= c_B^\top \cdot y_B + c_N^\top y_N \\ &= c_B^\top \cdot \left(A_B^{-1} \cdot b - A_B^{-1} A_N\, y_N\right) + c_N^\top\, y_N \\ &= c_B{}^\top \cdot \underbrace{A_B^{-1} \cdot b}_{=x_B} + \sum_{j \in N} \underbrace{\left(c_j - c_B^\top \cdot A_B^{-1} \cdot A_j\right)}_{=\bar{c}_j} y_j \\ &= \underbrace{c_B^\top \cdot x_B}_{=c^\top \cdot x} + \sum_{j \in N} \bar{c}_j\, y_j \geq c^\top \cdot x \end{aligned}$$

b

Assume by contradiction that $\bar{c}_j < 0$ for some $j \in N$.

Since x is non-degenerate, the jth basic direction is a feasible direction

and the cost can thus be decreased as $\bar{c}_j < 0$. □

Development of the Simplex Method

Assumption (for now): only *non-degenerate* basic feasible solutions

Let x be a basic feasible solution with $\bar{c}_j < 0$ for some $j \in N$.

Let d be the jth basic direction:

$$0 > \bar{c}_j = c^\top \cdot d$$

It is desirable to go to $y := x + \theta^* d$ with $\theta^* := \max\{\theta \mid x + \theta\, d \in P\}$.

Question: How to determine θ^*?

By construction of d, it holds that $A \cdot (x + \theta\, d) = b$ for all $\theta \in \mathbb{R}$, i.e.,

$$x + \theta\, d \in P \qquad \Longleftrightarrow \qquad x + \theta\, d \geq 0.$$

Case 1: $d \geq 0 \quad \Longrightarrow \quad x + \theta\, d \geq 0$ for all $\theta \geq 0 \quad \Longrightarrow \quad \theta^* = \infty$

Thus, the LP is unbounded.

Case 2: $d_k < 0$ for some $k \quad \Longrightarrow \quad \left(x_k + \theta\, d_k \geq 0 \quad \Longleftrightarrow \quad \theta \leq \dfrac{-x_k}{d_k} \right)$

Thus, $\theta^* = \min\limits_{k:\ d_k < 0} \dfrac{-x_k}{d_k} = \min\limits_{\substack{i=1,\dots,m \\ d_{B(i)} < 0}} \dfrac{-x_{B(i)}}{d_{B(i)}} > 0.$

Developement of the Simplex Method (Cont.)

Assumption (for now): only *non-degenerate* basic feasible solutions

Let x be a basic feasible solution with $\bar{c}_j < 0$ for some $j \neq B(1), \dots, B(m)$.

Let d be the jth basic direction:

$$0 > \bar{c}_j = c^\top \cdot d$$

It is desirable to go to $y := x + \theta^* \cdot d$ with $\theta^* := \max\{\theta \mid x + \theta \cdot d \in P\}$.

$$\theta^* = \min_{k:\ d_k < 0} \frac{-x_k}{d_k} = \min_{\substack{i=1,\dots,m \\ d_{B(i)} < 0}} \frac{-x_{B(i)}}{d_{B(i)}}$$

Let $\ell \in \{1, \dots, m\}$ with $\theta^* = \dfrac{-x_{B(\ell)}}{d_{B(\ell)}}$, then $y_j = \theta^*$ and $y_{B(\ell)} = 0$.

$\Longrightarrow$ x_j replaces $x_{B(\ell)}$ as a basic variable and we get a new basis matrix

$$A_{\bar{B}} = \left[A_{B(1)}, \dots, A_{B(\ell-1)}, A_j, A_{B(\ell+1)}, \dots, A_{B(m)}\right] = \left[A_{\bar{B}(1)}, \dots, A_{\bar{B}(m)}\right]$$

with $$\bar{B}(i) = \begin{cases} B(i) & \text{if } i \neq \ell, \\ j & \text{if } i = \ell. \end{cases}$$

Theorem 4.8 Let x be a non-degenerate basic feasible solution, $j \in N$ with $\bar{c}_j < 0$, d the jth basic direction, and $\theta^* := \max\{\theta \mid x + \theta d \in P\} < \infty$.

a $\theta^* = \min_{\substack{i=1,\dots,m \\ d_{B(i)}<0}} \frac{-x_{B(i)}}{d_{B(i)}} = \frac{-x_{B(\ell)}}{d_{B(\ell)}}$ for some $\ell \in \{1, \dots, m\}$.

Let $\bar{B}(i) := B(i)$ for $i \neq \ell$ and $\bar{B}(\ell) := j$.

b $A_{\bar{B}(1)}, \dots, A_{\bar{B}(m)}$ are linearly independent and $A_{\bar{B}}$ is a basis matrix.

c $y := x + \theta^* d$ is a basic feasible solution associated with $\bar{B}$ and $c^\top \cdot y < c^\top \cdot x$.

Proof: a we just calculated

b
$$\begin{aligned} A_B^{-1} A_{\bar{B}} &= A_B^{-1}[A_{B(1)}, \dots, A_{B(\ell-1)}, A_j, A_{B(\ell+1)}, \dots, A_{B(m)}] \\ &= [e_1, \dots, e_{\ell-1}, -d_B, e_{\ell+1}, \dots, e_m] \end{aligned}$$

Since $-d_{B(\ell)} > 0$, $A_{\bar{B}}$ has full rank, and $\bar{B}$ is a basis

c clear since $y_\ell = 0$, $y_j = \theta^*$, $\theta^* > 0$, and $\bar{c}_j < 0$

An Iteration of the Simplex Method

Given: basis B corresponding to basic feasible solution x

i Let $\bar{c}^\top := c^\top - c_B^\top \cdot A_B^{-1} \cdot A$. If $\bar{c} \geq 0$, then STOP; else choose j with $\bar{c}_j < 0$.

ii Let $u := A_B^{-1} \cdot A_j$. If $u \leq 0$, then STOP (optimal cost is $-\infty$).

iii Let $\theta^* := \min_{i\,:\,u_i>0} \frac{x_{B(i)}}{u_i} = \frac{x_{B(\ell)}}{u_\ell}$ for some $\ell \in \{1, \ldots, m\}$.

iv Form new basis by replacing $A_{B(\ell)}$ with A_j; corresponding basic feasible solution y is given by

$$y_j := \theta^* \qquad \text{and} \qquad y_{B(i)} = x_{B(i)} - \theta^* u_i \quad \text{for } i \neq \ell.$$

Remark: We say that the nonbasic variable x_j enters the basis and the basic variable $x_{B(\ell)}$ leaves the basis.

Correctness of the Simplex Method

Theorem 4.9 If every basic feasible solution is non-degenerate, the simplex method terminates after finitely many iterations in one of the following two states:

a. we have an optimal basis B and an associated basic feasible solution x which is optimal;
b. we have a vector d satisfying $A \cdot d = 0$, $d \geq 0$, and $c^\top \cdot d < 0$;
   the optimal cost is $-\infty$.

Proof sketch: The simplex method makes progress in every iteration.
Since there are only finitely many different basic feasible solutions,
it stops after a finite number of iterations. □

Introduction to

Linear and Combinatorial Optimization

4 Simplex Method

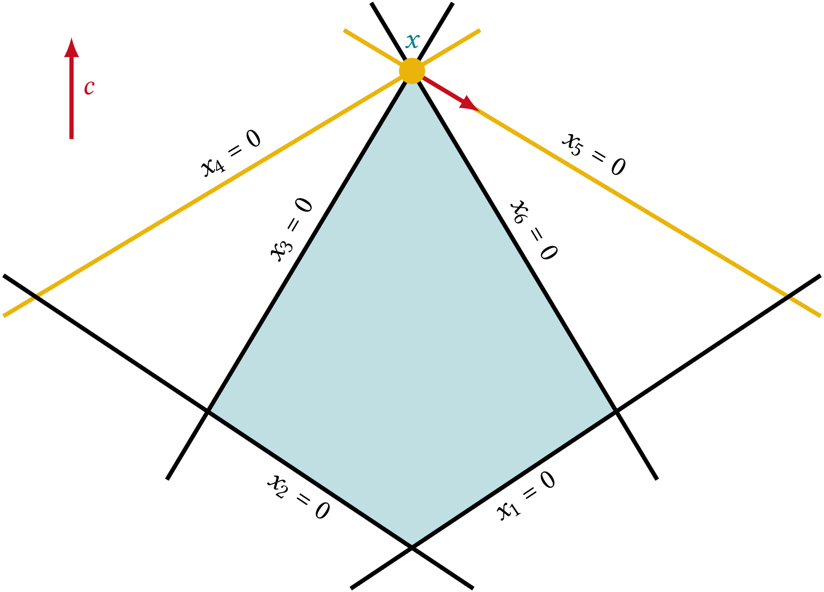
4.2 Degenerate Problems

Simplex Method for Degenerate Problems

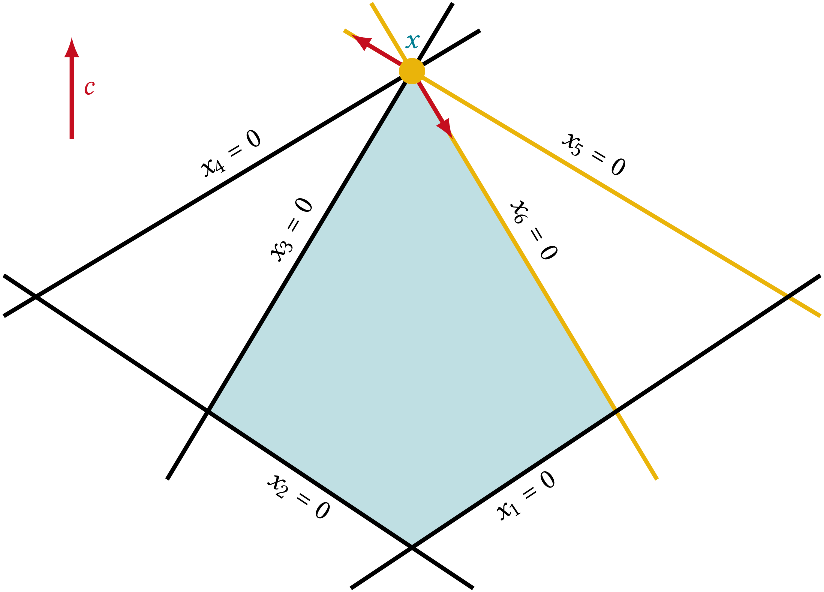
- An iteration of the simplex method can also be applied if x is a degenerate basic feasible solution.
- In this case it might happen that $\theta^* := \min_{i\,:\,u_i>0} \frac{x_{B(i)}}{u_i} = \frac{x_{B(\ell)}}{u_\ell} = 0$ if some basic variable $x_{B(\ell)}$ is zero and $d_{B(\ell)} = -u_\ell < 0$.
- Thus, $y = x + \theta^* d = x$ and the current basic feasible solution does not change.
- But replacing $A_{B(\ell)}$ with A_j still yields a new basis with associated basic feasible solution $y = x$.

Remark: Even if θ^* is positive, more than one of the original basic variables may become zero at the new point $x + \theta^* d$. Since only one of them leaves the basis, the new basic feasible solution y may be degenerate.

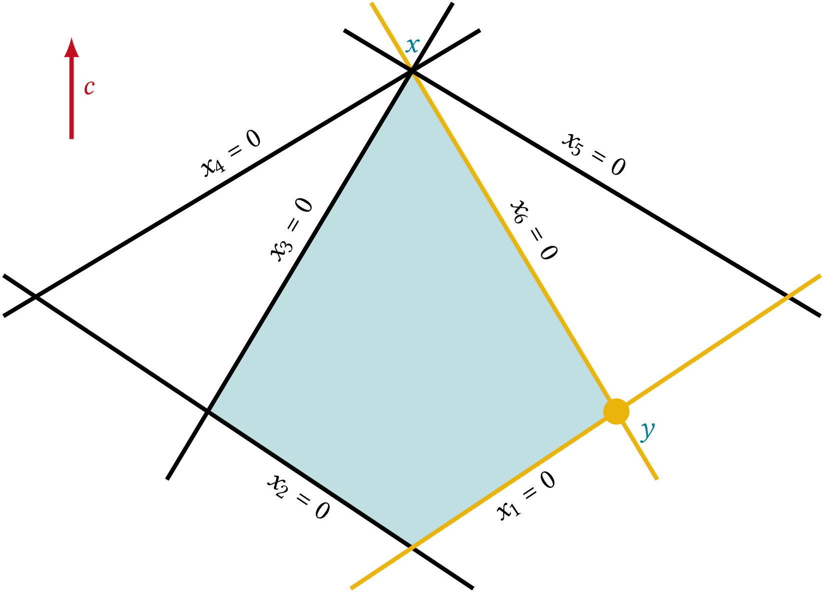
Example



Example



Example



Pivot Selection

Question: How to choose j with $\bar{c}_j < 0$ and ℓ with $\dfrac{x_{B(\ell)}}{u_\ell} = \min_{i\,:\,u_i>0} \dfrac{x_{B(i)}}{u_i}$ if several possible choices exist?

Attention: Choice of j is critical for overall behavior of simplex method.

Three popular choices are:

- smallest subscript rule: choose smallest j with $\bar{c}_j < 0$.
  (very simple; no need to compute entire vector $\bar{c}$; usually leads to many iterations)
- steepest descent rule: choose j such that $\bar{c}_j < 0$ is minimal.
  (relatively simple; commonly used for mid-size problems; does not necessarily yield the best neighboring solution)
- best improvement rule: choose j such that $\theta^* \bar{c}_j$ is minimal.
  (computationally expensive; used for large problems; usually leads to very few iterations)